

Network Structure and Activity in Boolean Networks

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with

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Boolean Networks (BN)

- ▶ A Boolean Network (BN) is a map of the form

$$F = (f_1, \dots, f_n): \{0, 1\}^n \longrightarrow \{0, 1\}^n .$$

- ▶ There are several types of stability concepts for dynamical systems. Examples:
 - Stability with respect to initial conditions
 - Attractor stability
 - Structural stability

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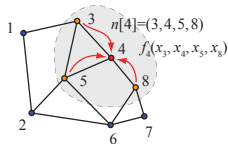
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- ▶ Here: Stability with respect to initial conditions. Specifically:

What is the probability that $F(x)$ and $F(x + e_i)$ are different?

Background, Terminology and Definitions – I

- ▶ Associated to F we have the *dependency graph of variables* denoted by X_F . It has vertex set $\{1, 2, \dots, n\}$ and edges all (i, j) for which f_j depend non-trivially on x_j .
- ▶ We let X denote the undirected version of X_F where loops are omitted.
- ▶ More definitions:
 - Ordered 1-neighborhood: $n[i]$
 - System state: $x = (x_1, \dots, x_n)$
 - Restricted vertex state: $x[i]$



Background, Terminology and Definitions – II

► Define the function $\alpha_{F,i}: K^n \rightarrow \{0, 1\}$ by

$$\alpha_{F,i}(x) = \mathbb{I}[F(x + e_i) \neq F(x)]$$

where \mathbb{I} is the indicator function and e_i is the i^{th} unit vector.

Definition

The *activity of F with respect to vertex i* is the expectation value of $\alpha_{F,i}$ using the uniform measure on K^n :

$$\bar{\alpha}_{F,i} = \mathbb{E}[\alpha_{F,i}] .$$

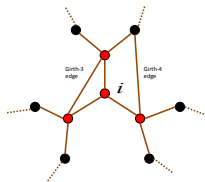
The *activity of F* is the vector

$$\bar{\alpha}_F = (\bar{\alpha}_{F,1}, \bar{\alpha}_{F,2}, \dots, \bar{\alpha}_{F,n}) .$$

Comments and Definitions

► Interpretation: For a randomly chosen state $x \in K^n$, the value $\bar{\alpha}_{F_i}$ may be regarded as the probability that perturbing x_i will cause $F(x + e_i) \neq F(x)$ to hold. This activity notion may naturally be regarded as a measure of sensitivity with respect to initial conditions.

The subgraph $X(i; 2)$ of X is induced by vertex i and its distance ≤ 2 neighbors.



► Vertices belonging to the closed 1-neighborhood $n[i]$ of i are marked red. *Type-3 edges* (relative to i) connect neighbors of i , while *type-4 edges* connect neighbors of $j \in n'[i]$ through a common neighbor different from i .

► Observation: $\bar{\alpha}_{F_i}(x)$ depends on the vertex states $X(i; 2)$ and the vertex functions f_j with $j \in n[i]$ (the set of vertices in the 1-neighborhood of i).

Related Work

- ▶ Shmulevich and Kauffman (2003) introduced a notion of activity for regular, random Boolean networks with a common vertex function f (see [3]). Here:

$$\alpha_{f,i} := \mathbb{E}\left[\frac{\partial f}{\partial x_i}\right]$$

- ▶ May be a reasonable measure for activity for the class of dynamical systems that they studied. However: does not incorporate the impact of network structure.
- ▶ Layne et. al. (2012) used the above definition to compute activity for canalizing functions [2].
- ▶ The activity notion introduced here incorporates network structure.

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- ▶ Layne et. al. (2012) used the above definition to compute activity for canalizing functions [2].
- ▶ The activity notion introduced here incorporates network structure.
- ▶ Why study activity? What are potential uses?

Preliminary Results

- ▶ We set $K = \{0, 1\}$, write N_i for the size of $X(i; 2)$, and $K(i) = K^{N_i}$ for the projection of K^n onto the set of vertex states associated to $X(i; 2)$.
- ▶ For $j \in n[i]$, define the sets $A_j(i) \subset K(i)$ and $A_j^0, A_j^1 \subset A_j(i)$ by

$$A_j(i) = \{x \in K(i) \mid F(x + e_i)_j \neq F(x)_j\} \quad \text{and} \quad A_j^m(i) = \{x \in A_j(i) \mid x_i = m\}, m = 0, 1.$$

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Proposition

For any Boolean network we have

$$\mathbb{P}\left(\bigcup_{j \in n[i]} A_j \mid x_i = 0\right) = \mathbb{P}\left(\bigcup_{j \in n[i]} A_j \mid x_i = 1\right).$$

Proof.

This follows by observing that there is a bijection $\phi_j: A_j^0 \rightarrow A_j^1$ since $x \in A_j^0$ if and only if $x + e_i \in A_j^1$ for all $j \in n[i]$.

Proposition

Let X be a graph and F a Boolean network map over X . The activity of F with respect to vertex i is

$$\bar{\alpha}_{F,i} = \mathbb{P}\left(\bigcup_{j \in n[i]} A_j \mid x_i = 0\right) = \mathbb{P}\left(\bigcup_{j \in n[i]} \bar{A}_j\right).$$

(Here $\bar{A}_j := A_j^0$)

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Proof.

$$\begin{aligned}\bar{\alpha}_{F,i} &= \mathbb{E}[\alpha_{F,i}] = \sum_{x \in K^n} \alpha_{F,i}(x) \mathbb{P}(x) = \sum_{x \in K(i)} \alpha_{F,i}(x) \mathbb{P}(x) = \mathbb{P}\left[\bigcup_{j \in n[i]} A_j\right] \\ &= \sum_{m=0,1} \mathbb{P}\left[\bigcup_{j \in n[i]} A_j \mid x_i = m\right] \mathbb{P}(x_i = m) = \mathbb{P}\left[\bigcup_{j \in n[i]} A_j \mid x_i = 0\right] = \mathbb{P}\left[\bigcup_{j \in n[i]} \bar{A}_j\right].\end{aligned}$$

(Here $\bar{A}_j := A_j^0$)

Examples

Recall that the standard Boolean *threshold function* $\tau_{k,n}: K^n \rightarrow K$ is defined by

$$\tau_{k,n}(x_1, \dots, x_n) = \begin{cases} 1, & \text{if } \sum_{j=1}^n x_j \geq k, \\ 0, & \text{otherwise,} \end{cases}$$

and that the logical nor-function $\text{nor}_m: K^m \rightarrow K$ is defined by

$$\text{nor}_m(x_1, \dots, x_m) = (1 + x_1) \cdots (1 + x_m),$$

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Proposition

If F is the BN map induced by the nor-function over K_n , then

$$\bar{\alpha}_{F,i} = \frac{1}{2^{n-1}},$$

and if F is induced by the k -threshold function over K_n , then

$$\bar{\alpha}_{F,i} = \binom{n-1}{k-1} / 2^{n-1}.$$

Proof.

For either choice of function, we have $\bar{A}_j = \bar{A}_k$ for all $j, k \in n[i]$.

For threshold functions, and for a state x with $x_i = 0$ to satisfy $\tau_k(x) \neq \tau_k(x + e_i)$, it is necessary and sufficient that x belong to Hamming class $k - 1$.

Since $x_i = 0$, it follows from Proposition 2 that $\mathbb{P}(\bar{A}_j(i)) = \frac{1}{2^{n-1}} |A_j| = \binom{n-1}{k-1} / 2^{n-1}$ as stated.

In the case of nor-functions we have $F(x + e_i)_i \neq F(x)_i$ precisely when $x_j = 0$ for all $j \in n(i)$ leading to $\mathbb{P}(\bar{A}_j) = \frac{1}{2^{n-1}}$.

Elementary Cellular Automata

Proposition (Threshold ECA)

The activity of k -threshold ECA for $n \geq 5$ is

$$\bar{\alpha}_{F,i} = \begin{cases} 0, & \text{if } k = 0 \text{ or } k > 3, \\ 1/2, & \text{if } k = 1 \text{ or } k = 3, \\ 7/8, & \text{if } k = 2. \end{cases}$$

For $n = 4$ and $k = 2$, the activity is $3/4$.

Proof (case $k = 1$).

We have

$$\bar{A}_{i-1} = \{(0, 0, 0, x_{i+1}, x_{i+2})\}, \quad \bar{A}_i = \{(x_{i-2}, 0, 0, 0, x_{i+2})\}, \quad \text{and} \quad \bar{A}_{i+1} = \{(x_{i-1}, x_{i-1}, 0, 0, 0)\}.$$

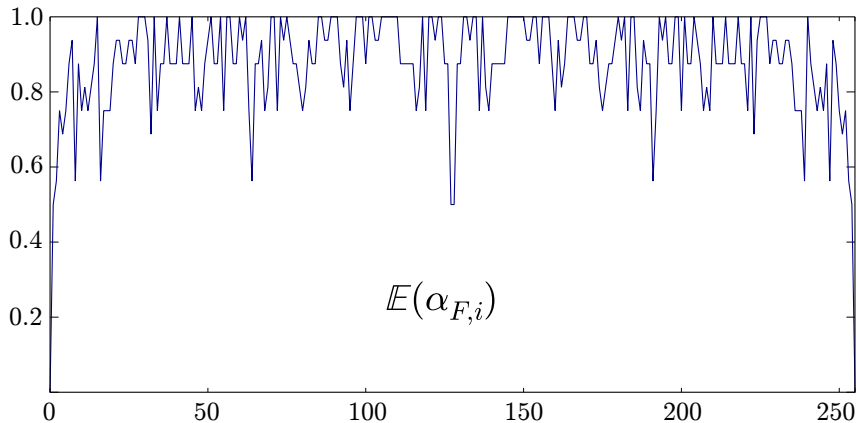
By the inclusion-exclusion principle:

$$\begin{aligned} |\bar{A}_{i-1} \cup \bar{A}_i \cup \bar{A}_{i+1}| &= |\bar{A}_{i-1}| + |\bar{A}_i| + |\bar{A}_{i+1}| \\ &\quad - |\bar{A}_{i-1} \cap \bar{A}_i| - |\bar{A}_{i-1} \cap \bar{A}_{i+1}| - |\bar{A}_i \cap \bar{A}_{i+1}| + |\bar{A}_{i-1} \cap \bar{A}_i \cap \bar{A}_{i+1}| \\ &= 3 \times 4 - 2 - 1 - 2 + 1 = 8 \end{aligned}$$

This yields $\bar{\alpha}_{F,i} = 8/2^4 = 1/2$.

We remark that if we instead use the `nor`-function, then $\bar{\alpha}_{F,i} = 1/2$ when $n \geq 5$.

Activity of Elementary Cellular Automata



More Elaborate Cases

- ▶ The evaluation of activity is more involved with dependencies/overlap among the A_j 's (except for at vertex i – of course).
- ▶ Will cover two cases here:
 - d -regular trees with threshold and nor-functions
 - Square lattices with nor-functions
- ▶ Additional cases covered in the conference paper.

d -regular trees

Proposition

Let X be a d -regular graph of girth ≥ 5 and F the GDS map over X induced by the nor-function. Then the activity of F with respect to i is given by

$$\bar{\alpha}_{F,i} = 1 - \left(1 - \frac{1}{2^d}\right)^d + \left(\frac{1}{2} - \frac{1}{2^d}\right)^d. \quad (1)$$

Note: let B denote the union of A_j 's with $j \neq i$. Then:

$$\mathbb{P}\left[\bigcup_{j \in n[i]} A_j\right] = \mathbb{P}(A_i \cup B) = \mathbb{P}(B) + \mathbb{P}(A_i \cap B^c) = 1 - \mathbb{P}(B^c) + \mathbb{P}(A_i \cap B^c),$$

where B^c denotes the complement of B .

Proof.

Conditioning on $x_i = 0$ and using independence we have

$$\mathbb{P}(\bar{B}^c) = \mathbb{P}\left(\bigcap_{j \in n(i)} \bar{A}_j^c\right) = \prod_{j \in n(i)} \mathbb{P}(\bar{A}_j) .$$

For a state x with $x_i = 0$ to be in \bar{A}_j all remaining d states of \bar{A}_j must be zero, leading to $\mathbb{P}(\bar{A}_j^c) = 1 - \frac{1}{2^d}$ and

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In order to calculate $\mathbb{P}(A_i \cap B^c)$ we note that $\mathbb{P}(A_i \cap B^c) = \mathbb{P}(A_i) \mathbb{P}(B^c | A_i)$ and again use independence to obtain

$$\mathbb{P}(\bar{B}^c | \bar{A}_i) = \prod_{j \in n(i)} \mathbb{P}(\bar{A}_j^c | \bar{A}_i) .$$

Note that $\mathbb{P}(\bar{A}_j | \bar{A}_i) = 1/2^{d-1}$ so that $\mathbb{P}(\bar{A}_j^c | \bar{A}_i) = 1 - \mathbb{P}(\bar{A}_j | \bar{A}_i) = 1 - 1/2^{d-1}$ which substituted back in gives

$$\mathbb{P}(\bar{B}^c | \bar{A}_i) = \left(1 - \frac{1}{2^{d-1}}\right)^d .$$

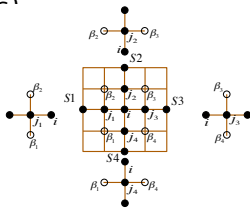
Noting that $\mathbb{P}(\bar{A}_i) = 1/2^d$ we obtain the third term in the proposition statement, finishing the proof

Proposition

If F is the GDS map with k -threshold vertex functions over a d -regular graph of girth ≥ 5 , then the activity of F with respect to vertex i is

$$\bar{\alpha}_{F,i} = 1 - \left[\frac{2^d - \binom{d}{k-1}}{2^d} \right]^d + \frac{\binom{d}{k-1}}{2^d} \left[\frac{2^{d-1} - \binom{d-1}{k-2}}{2^{d-1}} \right]^{k-1} \left[\frac{2^{d-1} - \binom{d-1}{k-1}}{2^{d-1}} \right]^{d-(k-1)} .$$

Square lattices (type-4 edge⁻¹)



Proposition

Let X be the 2-dimensional lattice (infinite or periodic boundaries) as above where every vertex has degree 4, and let F be the GDS map over X induced by nor-functions. Then the activity of F for any vertex i is $\bar{\alpha}_{F,i} = 1040/2^{12}$.

Proof.

Idea: use inclusion-exclusion and careful book-keeping twice to evaluate $\mathbb{P}(B)$ and $\mathbb{P}(\bar{A}_i \cap \bar{B}^c)$.

Remark

May of course obtain this through brute force evaluation: only 2^{13} state involved for $X(i; 2)$.




Questions, Acknowledgements and References

- Computational intractability: in the general case, determination of activity is intractable.
- Are there efficient and accurate methods for estimation?
- Activity only covers 1-step effects. Can something be done for long-term activity (e.g. with comparison of limit sets $\omega_F(x) \neq \omega_F(x + e_i)$)?
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